



Tutorial OmpSs

Agenda

9:00 – 10:30	Introduction to StarSs OmpSs syntax Simple examples Development methodology and infrastructure	90 min
10:30 – 11:00	Coffee break	30 min
11:10 – 12:30	Hands-on single node (I)	90 min
12:30 – 13:30	Lunch	60 min
13:30 – 15:00	Support for heterogeneous platforms Advanced examples	90 min
15:00 – 15:15	Short break	15 min
15:15 – 17:00	Hands-on (II)	105 min

BCS Barcelona Supercomputing Center
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Giving hints to the compiler: priorities

```

for (k = 0; k < nt; k++) {
    for (i = 0; i < k; i++) {
        #pragma omp task in([ts*ts]Ah[i*nt + k]) inout([ts*ts]Ah[k*nt + k]) \
        priority( (nt-i)+10 ) firstprivate (i, k, nt, ts)
        syrk_tile (Ah[i*nt + k], Ah[k*nt + k], ts, region)
    }
    // Diagonal Block factorization and panel permutations
    #pragma omp task inout([ts*ts]Ah[k*nt + k]) \
    priority( 100000 ) firstprivate (k, ts, nt)
    potr_tile(Ah[k*nt + k], ts, region)

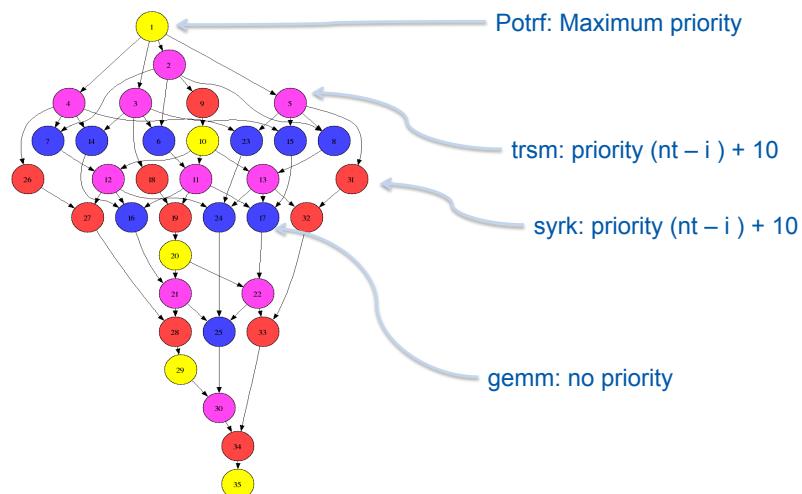
    // update trailing matrix
    for (i = k + 1; i < nt; i++) {
        for (j = 0; j < k; j++) {
            #pragma omp task in ([ts*ts]Ah[j*nt+i], [ts*ts]Ah[j*nt+k]) \
            inout ( [ts*ts]Ah[k*nt+i] ) firstprivate (i, j, k, ts, nt)
            gemm_tile (Ah[j*nt + i], Ah[j*nt + k], Ah[k*nt + i], ts, region)
        }
        #pragma omp task in([ts*ts]Ah[k*nt + k]) inout([ts*ts]Ah[k*nt + i]) \
        priority( (nt-i)+10 ) firstprivate (i, k, ts, nt)
        trsm_tile (Ah[k*nt + k], Ah[k*nt + i], ts, region)
    }
}
#pragma omp taskwait

```



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Giving hints to the compiler: priorities



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Giving hints to the compiler: priorities

« Two policies available:

- Priority

- Tasks are scheduled based on the assigned priority
- The priority is a number ≥ 0 . Given two tasks with priority A and B, where $A > B$, the task with priority A will be executed earlier than the one with B
- When a task T with priority A creates a task Tc that was given priority B by the user, the priority of Tc will be added to that of its parent. Thus, the priority of Tc will be $A + B$

- Smart Priority

- Similar to the Priority, but also propagates the priority to the immediate preceding tasks

« Using the schedulers:

- `export NX_ARGS= "--schedule-priority"`
- `export NX_ARGS = "--schedule-smartpriority"`

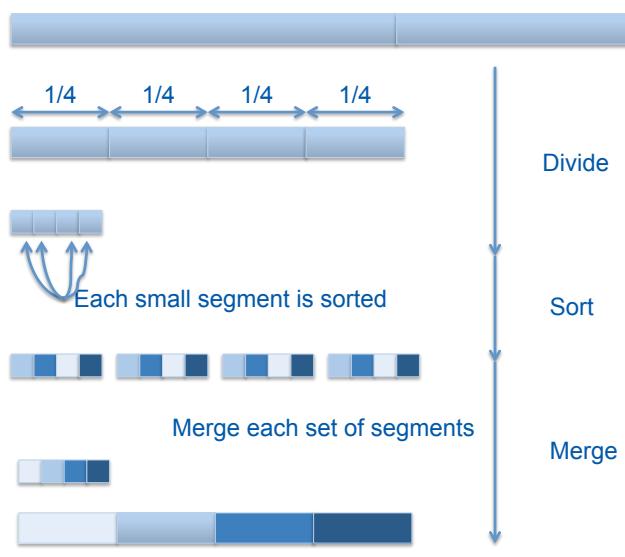
« Right now only available in bf, dbf schedulers



Accessing non-contiguous or partially overlapped regions

« Sorting arrays

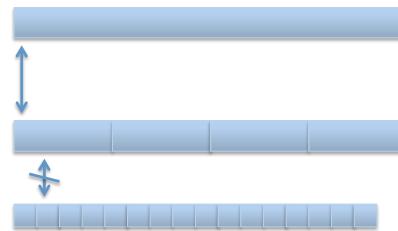
- Divide by $\frac{1}{4}$
- Sort
- Merge



Accessing non-contiguous or partially overlapped regions

« Why is the regions-aware dependences plug-in needed?

- Regular dependence checking uses first element as representative (size is not considered)
- Segment starting at address $A[i]$ with length $L/4$ will be considered the same as $A[i]$ with length L
- Dependences between $A[i]$ with length L and $A[i+L/4]$ with length $L/4$ will not be detected



« All these is fixed with the regions plugin

« Two different implementations:

- **NX_DEPS= regions**
- **NX_DEPS= perfect-regions**



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Accessing non-contiguous or partially overlapped regions

```
void multisort(long n, T data[n], T tmp[n]) {
    if (n >= MIN_SORT_SIZE*4L) {
        // Recursive decomposition
        #pragma omp task inout (data[0;n/4L]) firstprivate(n)
        multisort(n/4L, &data[0], &tmp[0]);
        #pragma omp task inout(data[n/4L;n/4L]) firstprivate(n)
        multisort(n/4L, &data[n/4L], &tmp[n/4L]);
        #pragma omp task inout (data[n/2L;n/4L]) firstprivate(n)
        multisort(n/4L, &data[n/2L], &tmp[n/2L]);
        #pragma omp task inout (data[3L*n/4L; n/4L]) firstprivate(n)
        multisort(n/4L, &data[3L*n/4L], &tmp[3L*n/4L]);

        #pragma omp task in (data[0;n/4L], data[n/4L;n/4L]) out (tmp[0; n/2L])\
        firstprivate(n)
        merge_rec(n/4L, &data[0], &data[n/4L], &tmp[0], 0, n/2L);
        #pragma omp task in (data[n/2L;n/4L], data[3L*n/4L; n/4L])\
        out (tmp[n/2L; n/2L]) firstprivate (n)
        merge_rec(n/4L, &data[n/2L], &data[3L*n/4L], &tmp[n/2L], 0, n/2L);

        #pragma omp task in (tmp[0; n/2L], tmp[n/2L; n/2L]) out (data[0; n]) firstprivate (n)
        merge_rec(n/2L, &tmp[0], &tmp[n/2L], &data[0], 0, n);
    }
    else basicsort(n, data);
}
```

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Accessing non-contiguous or partially overlapped regions

```

void merge_rec(long n, T left[n], T right[n], T result[n*2], long start, long length) {
    if (length < MIN_MERGE_SIZE*2L) {
        // Base case
        basicmerge(n, left, right, result, start, length);
    } else {
        // Recursive decomposition
        #pragma omp task in (left[0:n], right[0:n])\
        out (result[start:length/2]) firstprivate(n, start, length)
        merge_rec(n, left, right, result, start, length/2);
        #pragma omp task in (left[0:n], right[0:n])\
        out (result[start+length/2:length/2]) firstprivate(n, start, length)
        merge_rec(n, left, right, result, start + length/2, length/2);
    }
}

```



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Accessing non-contiguous or partially overlapped regions

```

T *data = malloc(N*sizeof(T));
T *tmp = malloc(N*sizeof(T));

posix_memalign ((void**)&data, N*sizeof(T), N*sizeof(T));
posix_memalign ((void**)&tmp, N*sizeof(T), N*sizeof(T));
. .
multisort(N, data, tmp);
#pragma omp taskwait

```

Current implementation requires alignment of data for efficient data-dependence management



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Using task versions

```
#pragma omp target device (smp) copy_deps
#pragma omp task input([NB][NB]A, [NB][NB]B) inout([NB][NB]C)
void matmul(double *A, double *B, double *C, unsigned long NB)
{
    int i, j, k, I;
    double tmp;
    for (i = 0; i < NB; i++) {
        I=i*NB;
        for (j = 0; j < NB; j++) {
            tmp=C[I+j];
            for (k = 0; k < NB; k++)
                tmp+=A[I+k]*B[k*NB+j];
            C[I+j]=tmp;
        }
    }
}

#pragma omp target device (smp) implements (matmul) copy_deps
#pragma omp task input([NB][NB]A, [NB][NB]B) inout([NB][NB]C)
void matmul_mkl(double *A, double *B, double *C, unsigned long NB)
{
    cblas_dgemm(CblasRowMajor, CblasNoTrans, CblasNoTrans, NB, NB, NB, 1.0,
    (double *)A, NB, (double *)B, NB, 1.0, (double *)C, NB);
}
```

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Using task versions

```
void compute(struct timeval *start, struct timeval *stop, unsigned long NB, unsigned long DIM, double *A[DIM][DIM], double *B[DIM][DIM], double *C[DIM][DIM])
{
    unsigned i, j, k;

    gettimeofday(start,NULL);

    for (i = 0; i < DIM; i++)
        for (j = 0; j < DIM; j++)
            for (k = 0; k < DIM; k++)
                matmul ((double *)A[i][k], (double *)B[k][j], (double *)C[i][j], NB);

    #pragma omp taskwait
    gettimeofday(stop,NULL);
}
```

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Using task versions

« Use of specific scheduling:

- export NX_SCHEDULE=versioning

« Tries each version a given number of times and automatically will choose the best version



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Using socket aware scheduling

- « Assign top level tasks (depth 1) to a NUMA node set by the user before task creation
 - nested tasks will run in the same node as their parent.
- « nanos_current_socket API function must be called before instantiation of tasks to set the NUMA node the task will be assigned to.
- « Queues sorted by priority with as many queues as NUMA nodes specified (see num-sockets parameter).



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Using socket aware scheduling

Example: stream

```
#pragma omp task in ([bs]a, [bs]b) out ([bs]c)
void add_task (double *a, double *b, double *c, int bs)
{
    int j;
    for (j=0; j < BSIZE; j++)
        c[j] = a[j]+b[j];
}

void tuned_STREAM_Add()
{
    int j;
    for (j=0; j<N; j+=BSIZE){
        nanos_current_socket( ( j/((int)BSIZE) ) % 2 );
        add_task(&a[j], &b[j], &c[j], BSIZE);
    }
}
```



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Using socket aware scheduling

Usage:

- export NX_SCHEDULE=socket

If using less than 6 threads:

- export NX_ARGS="--binding-stride 6"



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